

Application: Constrained Optimization (cont'd)

Recall: a quadratic form is

$$q(x_1, \dots, x_n) = \sum a_{ij} x_i x_j \iff q(x) = x^T S x, \quad S \text{ symmetric}$$

Eg: $q(x_1, x_2, x_3) = a_{11}x_1^2 + a_{22}x_2^2 + a_{33}x_3^2 + a_{12}x_1x_2 + a_{13}x_1x_3 + a_{23}x_2x_3$

then $q(x) = x^T S x$ for $S = \frac{1}{2} \begin{pmatrix} 2a_{11} & a_{12} & a_{13} \\ a_{12} & 2a_{22} & a_{23} \\ a_{13} & a_{23} & 2a_{33} \end{pmatrix}$

Quadratic Optimization Problem:

Find min/max values of $q(x)$ subject to $\|x\|=1$.

Procedure for Solving a Quadratic Optimization Problem:

(1) Write $q(x) = x^T S x$ for S symmetric

(2) Orthogonally diagonalize S : *orthonormal eigenbasis*

$$S = Q D Q^T \quad Q = \begin{pmatrix} | & & | \\ v_1 & \dots & v_n \\ | & & | \end{pmatrix} \text{ orthogonal, } D = \begin{pmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_n \end{pmatrix}$$

Order so that $\lambda_1 \geq \dots \geq \lambda_n$.

→ Set $y = Q^T x$. Then

$$q(x) = x^T Q D Q^T x = (Q^T x) D (Q^T x) = y^T D y \\ = \lambda_1 y_1^2 + \dots + \lambda_n y_n^2$$

q is diagonal in the y -coordinates.

(3) The maximum value of q is the largest eigenvalue λ_1 :

$$q(\omega_1) = \lambda_1 \quad \text{b/c } Q^T \omega_1 = \begin{pmatrix} 1 \\ \vdots \\ 0 \end{pmatrix} \text{ i.e. } y_1 = 1, y_2 = \dots = y_n = 0$$

The minimum value of q is the smallest eigenvalue λ_n :

$$q(\omega_n) = \lambda_n \quad \text{b/c } Q^T \omega_n = \begin{pmatrix} 0 \\ \vdots \\ 1 \end{pmatrix} \text{ i.e. } y_1 = \dots = y_{n-1} = 0, y_n = 1$$

Eg: (from last time) $q(x_1, x_2) = \frac{5}{2}(x_1^2 + x_2^2) - x_1 x_2$
 $q(x) = x^T S x$ $S = \frac{1}{2} \begin{pmatrix} 5 & -1 \\ -1 & 5 \end{pmatrix} = Q^T D Q$
 $Q = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ -1 & 1 \end{pmatrix}$ $D = \begin{pmatrix} 3 & 0 \\ 0 & 2 \end{pmatrix}$
 \Rightarrow max is $q(1/\sqrt{2}, -1/\sqrt{2}) = 3$
 min is $q(1/\sqrt{2}, 1/\sqrt{2}) = 2$

Ellipsoids

Suppose all $\lambda_i > 0$. Then

$q(x) = \lambda_1 y_1^2 + \dots + \lambda_n y_n^2 = 1$ is an **ellipsoid**.

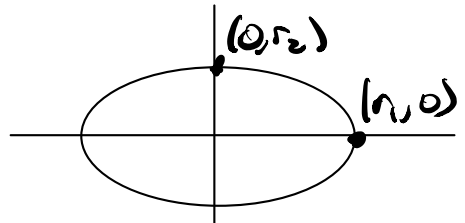
If you take the $(n-1)$ -sphere

$$x_1^2 + \dots + x_n^2 = 1$$

and change coords $y_i = \sqrt{\lambda_i} x_i$.

When $n=2$ you get an **ellipse**:

$$\left(\frac{x_1}{r_1}\right)^2 + \left(\frac{x_2}{r_2}\right)^2 = 1$$

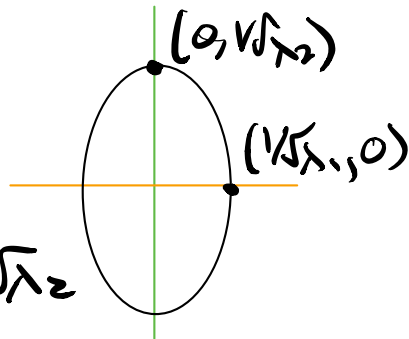


When you diagonalize $q(x_1, x_2)$, $y = Q^T x \Leftrightarrow x = Q y$

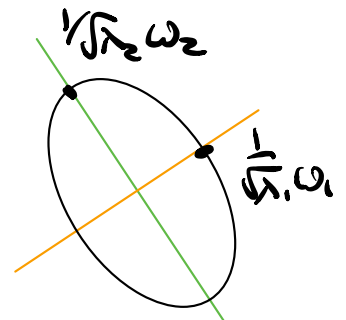
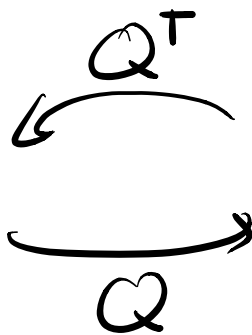
$$\lambda_1 \geq \lambda_2$$

$$\frac{1}{\sqrt{\lambda_1}} \leq \frac{1}{\sqrt{\lambda_2}}$$

$$r_1 = 1/\sqrt{\lambda_1} \quad r_2 = 1/\sqrt{\lambda_2}$$



$$\lambda_1 y_1^2 + \lambda_2 y_2^2 = 1$$



$$q(x_1, x_2) = 1$$

Relation to Quadratic Optimization:

Maximize $\|x\|$ subject to $q(x) = 1 \Rightarrow x = 1/\sqrt{\lambda_1} \omega_1$

$\Rightarrow q(\omega_1) = q(x/\|x\|) = 1/\|x\|^2 q(x) = \lambda_1$ maximized.

Maximize $\|x\|$ subject to $q(x)=1 \rightarrow x = \sqrt{\lambda_2} \omega_2$
 $\Rightarrow q(\omega_2) = q(x/\|x\|) = 1/\|x\|^2 q(x) = \lambda_2$ minimized.

Positive-Definite matrices

The condition "all $\lambda_i > 0$ " is special - and very common.

Def: A symmetric matrix is

- **positive-definite** if all eigenvalues are > 0
- **positive-semidefinite** if all eigenvalues are ≥ 0
- Likewise for **negative-definite** & **negative-semidefinite**.
- **Indefinite** otherwise.

+ve-def	+ve-semidet	-ve-def	indefinite
$Q \begin{pmatrix} 3 & 0 \\ 0 & 2 \end{pmatrix} Q^T$	$Q \begin{pmatrix} 3 & 0 \\ 0 & 0 \end{pmatrix} Q^T$	$Q \begin{pmatrix} -2 & 0 \\ 0 & -3 \end{pmatrix} Q^T$	$Q \begin{pmatrix} 3 & 0 \\ 0 & -2 \end{pmatrix} Q^T$

Positive-definiteness is an important condition! Want to check **without computing eigenvalues!**

Conditions for S to be Positive-Definite:

The following are equivalent for a symmetric matrix S :

- (1) All eigenvalues are > 0 +ve-def
- (2) $x^T S x > 0$ for all $x \neq 0$ +ve-energy
- (3) All n upper-left determinants are > 0 .

$$S = \begin{pmatrix} S_{11} & \dots & \dots \\ \vdots & \ddots & \vdots \\ \dots & \dots & \dots \end{pmatrix} \quad \det(S) > 0$$

etc.

(4) $S = A^T A$ for a matrix A with full col. rank.

(5) $S = LU$ (no row swaps) and diagonal entries of U are > 0 . ← fastest! (elimination problem)

Eg: $S = \begin{pmatrix} 7 & 2 & 0 \\ 2 & 6 & 2 \\ 0 & 2 & 5 \end{pmatrix}$ $p(\lambda) = -(\lambda-3)(\lambda-6)(\lambda-9)$

(1) Eigenvalues are $3, 6, 9 > 0$.

(2) $x^T S x = 9y_1^2 + 6y_2^2 + 3y_3^2 > 0$ $S = Q D Q^T$ $y = Q^T x$
1x1 matrix whenever $x \neq 0$ ($\Rightarrow Q^T x = y \neq 0$)

(3) $\det(7) = 7 > 0$ $\det \begin{pmatrix} 7 & 2 \\ 2 & 6 \end{pmatrix} = 38 > 0$ $\det(S) = 3 \cdot 6 \cdot 9 > 0$

(4) Cholesky: next

(5) $\begin{pmatrix} 7 & 2 & 0 \\ 2 & 6 & 2 \\ 0 & 2 & 5 \end{pmatrix} \rightsquigarrow \begin{pmatrix} 7 & 2 & 0 \\ 0 & 38/7 & 2 \\ 0 & 2 & 5 \end{pmatrix} \longrightarrow \begin{pmatrix} 7 & 2 & 0 \\ 0 & 38/7 & 2 \\ 0 & 0 & 4\frac{10}{38} \end{pmatrix}$
 U \uparrow diag > 0

Whence these conditions?

(1) is the definition

(2) $x^T S x = \lambda_1 y_1^2 + \dots + \lambda_n y_n^2$. All $\lambda_i > 0 \Rightarrow$
this is > 0 when $x \neq 0$ ($\Rightarrow y = Q^T x \neq 0$)

(3) Determinants: magic

(4) If $S = A^T A$ for A with full col rank:

$$x^T S x = x^T A^T A x = (Ax)^T (Ax) = (Ax) \cdot (Ax) = \|Ax\|^2$$

If A has FCR then $x \neq 0 \Rightarrow Ax \neq 0$.

(5) From $S = L D L^T$: next.

LDL^T and Cholesky.

Basically LU for +ve-def symmetric matrices.

Thm: A positive-definite symmetric matrix S can be uniquely decomposed as
 $S = LDL^T$ and $S = L_1 L_1^T$ ← Cholesky

where:

D : diagonal w/ positive diagonal entries

L : lower-unitriangular

L_1 : lower-triangular with positive diagonal entries.

Pf: [supplement]

NB: L^T has full col rank (it's invertible) and
 $S = A^T A$ for $A = L^T$. So this justifies condition 3.

NB: If $D = \begin{pmatrix} d_1 & & 0 \\ & \ddots & \\ 0 & & d_n \end{pmatrix}$ set $\sqrt{D} = \begin{pmatrix} \sqrt{d_1} & & 0 \\ & \ddots & \\ 0 & & \sqrt{d_n} \end{pmatrix}$

Define $L_1 = L\sqrt{D}$ (scales cols of L by $\sqrt{d_1}, \dots, \sqrt{d_n}$)

Then $L_1 L_1^T = L\sqrt{D}\sqrt{D}^T L^T = LDL^T = S$.

So LDL^T & $L_1 L_1^T$ are interchangeable.

We'll focus on LDL^T .

NB: If $S = LDL^T$ set $U = DL^T$ (scales rows of L^T by entries of D). Then $S = LU$ is the LU decomposition!

Procedure to compute $S = LDL^T$: S symmetric, +ve-def

(1) Find $S = LU$

Magic: never need row swaps!

(2) $L = L$ $D =$ diagonal entries of U : $S = LDL^T$

Magic: Diagonal entries of U are > 0 .

Magic: $DL^T = U$ i.e. $D^{-1}U = L^T$.

Why $LDL^T / L, L^T$?

- Solve $Sx = b$ by forward/back substitution (same reason you want LU)
- Can compute LDL^T about twice as fast as LU , & with half the memory.

LU : $\frac{2}{3}n^3$ flops

LDL^T : $\frac{1}{3}n^3$ flops

[supplement]