LDL' & Cholesky

This anounts to an LU decomposition of a positivedefinite, symmetric matrix that's 1x as fast to compute!

thm: A positive-definite symmetric matrix S can be uniquely decomposed as S=LDLT and S=LLT - Cholesky where: D: dragonal w/positive dragonal entries L: lower-unitriangular Li: lover-trangular with positive diagonal entries. Proof : [supplement] NB: Any such L, has full column rounts so S=LiLT is necessarily positive-definite & symmetric (last time). NB: Let U=DLT. (scales the rows of LT by the dragonal entries of D) Then U is upper-D with positive diagonal entries => in REF, so S=LU is the LU decomposition! This tells us how to compute an LDLT decomposition.

Procedure to compute S=LDL^T: Let S be a symmetric matrix. (1) Compute the LU decomposition S=LU. -> If you have to do a row swap then stop: Sis not positive-definite. -If the diagonal entries of U are not all positive then stop: Sis not positive-definite. (2) let D= the matrix of diagonal entries of U (set the off-diagonal entries = 0). Then $S = LDL^{T}$.

NB: An LDL⁺ decomposition can be computed in $\gamma_3^+ n^3$ flops (as opposed to 2/3 n³ for LU). This requires a slightly more dever absorption. See the supplement - it's also foster by hand!

NB: This is still on LU decomposition - lets you solve Sx=b quickly.

NB: S=QDQ³ and S=LDL⁷ are both "diagonalizations" in the serve of quadratic forms (later).

Find the LDLT decomposition of

$$S = \begin{pmatrix} 2 & 4 & -2 \\ 4 & q & -1 \\ -2 & -1 & 14 \end{pmatrix}$$
2-column L U
nothod:

$$\begin{pmatrix} 1 \\ 2 \\ -2 & -1 & 14 \end{pmatrix}$$

$$R_{1} = 2R_{1} \begin{pmatrix} 1 \\ 2 \\ -1 \end{pmatrix} \begin{pmatrix} 2 & 4 & -2 \\ 4 & q & -1 \\ -2 & -1 & 14 \end{pmatrix}$$

$$R_{2} = 2R_{1} \begin{pmatrix} 1 \\ 2 \\ -1 \end{pmatrix} \begin{pmatrix} 2 & 4 & -2 \\ 0 & 1 & 3 \\ 0 & 3 & 12 \end{pmatrix}$$

$$R_{3} = 3R_{2} \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 3 & 1 \end{pmatrix} \begin{pmatrix} 2 & 4 & -2 \\ 0 & 1 & 3 \\ 0 & 3 & 12 \end{pmatrix}$$

$$R_{3} = 3R_{2} \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 3 & 1 \end{pmatrix} D = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 3 \\ 0 & 0 & 3 \end{pmatrix}$$

$$S_{2} = LDLT \text{ for}$$

$$L = \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 3 & 1 \end{pmatrix} D = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 3 \\ 0 & 0 & 3 \end{pmatrix}$$

$$Check = DLT = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 3 \\ 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 2 & 4 & -2 \\ 0 & 1 & 3 \\ 0 & 0 & 3 \end{pmatrix} = (M_{1} + M_{2} + M_{2}$$

$$\begin{aligned} IF \quad D = \begin{pmatrix} d_{1} & 0 \\ 0 & d_{n} \end{pmatrix} \quad \text{set} \quad JD = \begin{pmatrix} VA_{1} & 0 \\ 0 & JA_{n} \end{pmatrix} \\ \text{Then} \quad JD \cdot JD = D \quad \text{and} \quad JD^{T} = JD, \text{ so} \\ \text{LDL}^{T} = L \cdot JD \cdot JD \cdot JT = (LJD)(LJD)^{T} \\ \text{So} \quad \text{just set} \\ L_{1} = LJD \implies S = L_{1}L_{1}^{T} \end{aligned}$$

$$\begin{aligned} \text{Strang:} & \text{"S = ATA} \quad \text{is how a positive-definite symmetric} \\ \text{medrix is put tagether.} \\ \text{S = L_{1}L_{1}^{T} \quad \text{is haw you pull it apart"} \end{aligned}$$

$$\begin{aligned} & \left(\begin{array}{c} 2 & 4 & -2 \\ -2 & -1 & 14 \end{array} \right) = L_{1}L_{1}^{T} \quad \text{for} \\ L_{1} = \begin{pmatrix} 1 & 0 & 0 \\ -3 & 1 & 0 \end{pmatrix} \begin{pmatrix} S_{2} & 0 & 0 \\ 0 & J_{2} & 0 \end{pmatrix} = \begin{pmatrix} J_{2} & 0 & 0 \\ -J_{2} & 3 & J_{3} \end{pmatrix} \end{aligned}$$

IF Sis positive-definite then S=LDUT

where D is diagonal with positive diagonal entries.

Cholesky from LDLT:

Quadratic Optimization This is an important application of the spectral theorem and positive-definiteness. Also, SVD+QO+e-stats=PCA.

It is the simplest case of quadratic programming, which is a big subfield of optimization. (So is least squares.)

For an example application, see the Wikipedia page for support-vector machine, an important tool in machine learning that reduces to a quadratic optimization problem. (There are tons of other applications.)

Def: An optimization problem means finding extremal values (minimum & maximum) of a function $f(x_1,...,x_n)$ subject to some constraint on $(x_{1,...,x_n})$.

In quadratic optimization, we consider quadratic functions. Def: A quadratic form in a variables is a function $q(x_1,...,x_n) = sum of terms of the form as x_i X_i$

Eq: $q(x_1, x_2) = \frac{1}{2}x_1^2 + \frac{1}{2}x_2^2 - x_1x_2$ Non eq: $q(x_1, x_2) = x_1^2 + x_2^2 + x_1 + x_2$ is not a quadratic form: x_1, x_2 are linear ferms.

NB: Thinking of
$$x = (x_0, ..., x_n)$$
 as a vector,
 $q(cx) = q(cx_1, ..., cx_n) = \sum_{i=1}^{n} a_{ij} (cx_i)(cx_j)$
 $= \sum_{i=1}^{n} c^2 a_{ij} x_i x_j = c^2 q(x)$
 $q(cx) = c^2 q(x)$

$$E_{g}: q(x_{1}, x_{2}) = 3x_{1}^{2} - 2x_{2}^{2}$$

$$\begin{aligned} & \text{Maximum}^{i} \\ & q(x_{ij}, x_{2}) = 3x_{1}^{2} - 2x_{2}^{2} \leq 3x_{1}^{2} + 3x_{2}^{2} \\ & = 3(x_{1}^{2} + x_{2}^{2}) = 3||x||^{2} = 3 \\ & \text{So the maximum value } 3 = 3; \text{ it is achieved} \\ & \text{at } (x_{ij}, x_{2}) = \pm (1, 0) : q(\pm 1, 0) = 3. \end{aligned}$$

Minimum

$$q(x_{v}x_{v}) = 3x_{i}^{2} - 2x_{i}^{2} = -2x_{i}^{2} - 2x_{i}^{2}$$
$$= -2(x_{i}^{2} + x_{i}^{2}) = -2||x_{i}|^{2} = -2$$

So the minimum value
$$\overline{3} - 2$$
; it is achieved
at $(x_i, x_i) = \pm (0, 1)$; $q(0, \pm 1) = -2$.

This example is easy because $q(x_1, x_2) = 3x_1^2 - 2x_2^2$ involves only squares of the coordinates: there is no cross-ferm XiXz

Strategy: To solve a quadratic aptomization problem, we want to diagonalize it to get n'd of the cross terms. To do this, we use symmetric matrices! Fact: Every guadratic form can be written $q(x) = x^T S x$ for a symmetric matrix S. Eq: $S = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 2 & 5 & c \end{pmatrix}$ $x_{T} X^{T} S x = (x_{1} \times x_{2} \times x_{3}) \begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{pmatrix} \begin{pmatrix} X_{1} \\ X_{2} \\ X_{3} \end{pmatrix}$ $= (X_{1} X_{2} X_{2}) \begin{pmatrix} X_{1} + 2 \times 2 + 3 \times 3 \\ 2 \times 1 + 4 \times 2 + 5 \times 3 \\ 3 \times 1 + 5 \times 2 + 5 \times 3 \end{pmatrix}$ $= \chi^{2} + 2\chi_{1}\chi_{2} + 3\chi_{1}\chi_{3}$ + 2x, x, + 4x2 + 5x, x3 + 3x5x1 + 5x3x1 + 6x52 = X2+4x2+6x3+4x1x2+6x1x3+10x2x3 NB: The (1,2) and (2,1) entries contribute to the XXX coefficient.

Given qs have to get S?
The xi² coefficients qo on the diagonals and
holf of the xix coefficient goes in the (ij) and
(i,i) arthes.

$$q(x_1, x_2, x_3) = a_1x_1^2 + a_{13}x_1^2 + a_{23}x_3^2 + a_{12}x_1x_2 + a_{23}x_3x_3 + a_{23}x_3 + a_{$$

Let
$$x = Qy$$
: this is a charge of variables
 $q(x) = q(Qy) = (Qy)^T Q D Q^T(Q_3)$
 $= y^T Q Q D Q^T Q_3 = y^T D y$
This is now disgonal!
NB: Q is a theorem \Rightarrow $||x|| = ||Qy|| = ||y||$
So $||x|| = 1 \Rightarrow ||y|| = 1$
Eq: Find the minimum & nextmum of
 $q(x_1, x_2) = \frac{1}{3}x_1^2 + \frac{1}{3}x_2^2 - 5x x_3 = constant
 $q(x_1) = x^T \begin{pmatrix} 1/3 & -5/2 \\ -5/2 & 1/2 \end{pmatrix} \times -9 = \frac{1}{2} \begin{pmatrix} 1 & -5 \\ -5 & 1 \end{pmatrix}$
Orthogonally diagonalize: $S = Q D Q^T$ for
 $Q = \frac{1}{32} \begin{pmatrix} -1 & 1 \\ 1 & 1 \end{pmatrix} D = \begin{pmatrix} 3 & 0 \\ 0 & -2 \end{pmatrix}$
Set $x = Qy$:
 $\begin{pmatrix} x_1 \\ y_2 \\ x_2 \end{pmatrix} = \frac{1}{32} \begin{pmatrix} -1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \frac{1}{32} \begin{pmatrix} -y_1 + y_2 \\ y_1 + y_2 \end{pmatrix}$
 $\begin{cases} x_1 = \frac{1}{32} (-y_1 + y_2) \\ x_2 = \frac{1}{32} (y_1 + y_2) \end{cases}$
Then $q(x) = y^T \begin{pmatrix} 3 & 0 \\ 0 & -2 \end{pmatrix} = 3y_1^2 - 2y_2^2$.$

Check:

$$q(x) = q\left(\frac{1}{2}\left(-\frac{1}{3}, \frac{1}{3}\right), \frac{1}{2}\left(\frac{1}{3}, \frac{1}{3}\right)\right)$$

$$= \frac{1}{2}\cdot\frac{1}{2}\left(-\frac{1}{3}, \frac{1}{3}\right), \frac{1}{2} + \frac{1}{2}\cdot\frac{1}{2}\left(\frac{1}{3}, \frac{1}{3}\right), \frac{1}{2} + \frac{1}{2}\cdot\frac{1}{3}\left(\frac{1}{3}, \frac{1}{3}\right), \frac{1}{2} + \frac{1}{2}\cdot\frac{1}{3}\left(\frac{1}{3}, \frac{1}{3}\right), \frac{1}{2} + \frac{1}{2}\cdot\frac{1}{3}\left(\frac{1}{3}, \frac{1}{3}\right), \frac{1}{2} + \frac{1}{4}\cdot\frac{1}{3}\right), \frac{1}{2}\cdot\frac{1}{2}\cdot\frac{1}{3}\cdot\frac{1$$

Quadratic Optimization= To find the minimum/maximum of a quadratic form q(x) subject to ||x||=1: (1) Write q(x)=xtSx for a symmetric matrix S (2) Orthogonally diagonalize S=QDQT for $Q = \begin{pmatrix} 1 & 1 \\ u_1 & \dots & u_n \end{pmatrix} \qquad D = \begin{pmatrix} \lambda_1 & 0 \\ 0 & \ddots & \lambda_n \end{pmatrix}$ eigenvæctors eigenvalues Order the eigenvalues so $\lambda_1 \ge \cdots \ge \lambda_n$ (3) The maximum value of g(x) is the largest eigenvalue λ_i . It is achieved for x = any unit λ_1 -eigenvector The minimum value of q(x) is the smallest ergenvalue Nn. It is achieved for x = any unit In-eigenvector. NB: If GM(ni)=1 then the only unit 2;-eigenvectors are ± vi. (only 2 unit vectors are on any line) NB: x=Qy diagonalizes q: ui -ui $q(x) = \lambda_i y_i^2 + \dots + \lambda_n y_n^2$