

# LDL<sup>T</sup> & Cholesky

This amounts to an LU decomposition of a **positive-definite, symmetric** matrix that's 2x as fast to compute!

**Thm:** A positive-definite symmetric matrix  $S$  can be uniquely decomposed as  
 $S = LDL^T$  and  $S = L_1 L_1^T \leftarrow \text{Cholesky}$

where:

$D$ : diagonal w/ positive diagonal entries

$L$ : lower-**unit**triangular

$L_1$ : lower-triangular with positive diagonal entries.

**Proof:** [supplement]

**NB:** Any such  $L_1$  has full column rank so  $S = L_1 L_1^T$  is necessarily positive-definite & symmetric (last time).

**NB:** Let  $U = DL^T$ .

(scales the rows of  $L^T$  by the diagonal entries of  $D$ )

then  $U$  is upper- $\Delta$  with positive diagonal entries

$\Rightarrow$  in REF, so  $S = LU$  is the LU decomposition!

This tells us how to compute an LDL<sup>T</sup> decomposition.

## Procedure to compute $S = LDL^T$ :

Let  $S$  be a symmetric matrix.

(1) Compute the LU decomposition  $S = LU$ .

→ If you have to do a row swap then **stop**:  
 $S$  is not positive-definite.

→ If the diagonal entries of  $U$  are not all positive then **stop**:  $S$  is not positive-definite.

(2) Let  $D$  = the matrix of diagonal entries of  $U$   
(set the off-diagonal entries = 0). Then  
 $S = LDL^T$ .

**NB:** An  $LDL^T$  decomposition can be computed in  $\sim \frac{1}{3}n^3$  flops (as opposed to  $\frac{2}{3}n^3$  for  $LU$ ). This requires a slightly more clever algorithm. See the **supplement** — it's also faster by hand!

**NB:** This is still an LU decomposition — lets you solve  $Sx = b$  quickly.

**NB:**  $S = QDQ^T$  and  $S = LDL^T$  are both "diagonalizations" in the sense of quadratic forms (later).

Eg: Find the LDL<sup>T</sup> decomposition of

$$S = \begin{pmatrix} 2 & 4 & -2 \\ 4 & 9 & -1 \\ -2 & -1 & 14 \end{pmatrix}$$

2-column  
method:

L

U

$$\begin{pmatrix} & & \\ & & \\ & & \end{pmatrix}$$

$$\begin{pmatrix} 2 & 4 & -2 \\ 4 & 9 & -1 \\ -2 & -1 & 14 \end{pmatrix}$$

$$\begin{array}{l} R_2 \leftarrow 2R_1 \\ R_3 \leftarrow R_1 \end{array}$$

$$\begin{pmatrix} 1 & & \\ 2 & & \\ -1 & & \end{pmatrix}$$

$$\begin{pmatrix} 2 & 4 & -2 \\ 0 & 1 & 3 \\ 0 & 3 & 12 \end{pmatrix}$$

$$R_3 \leftarrow 3R_2$$

$$\begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 3 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 2 & 4 & -2 \\ 0 & 1 & 3 \\ 0 & 0 & 3 \end{pmatrix}$$

So  $S = LDL^T$  for

$$L = \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 3 & 1 \end{pmatrix}$$

$$D = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 3 \end{pmatrix}$$

Check:

$$DL^T = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 3 \end{pmatrix} \begin{pmatrix} 1 & 2 & -1 \\ 0 & 1 & 3 \\ 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 2 & 4 & -2 \\ 0 & 1 & 3 \\ 0 & 0 & 3 \end{pmatrix} = U \quad \checkmark$$

## Cholesky from LDL<sup>T</sup>:

If  $S$  is positive-definite then  $S = LDL^T$  where  $D$  is diagonal with **positive** diagonal entries.

If  $D = \begin{pmatrix} d_1 & & 0 \\ & \ddots & \\ 0 & & d_n \end{pmatrix}$  set  $\sqrt{D} = \begin{pmatrix} \sqrt{d_1} & & 0 \\ & \ddots & \\ 0 & & \sqrt{d_n} \end{pmatrix}$

Then  $\sqrt{D} \cdot \sqrt{D} = D$  and  $\sqrt{D}^T = \sqrt{D}$ , so

$$LDL^T = L\sqrt{D}\sqrt{D}L^T = (L\sqrt{D})(L\sqrt{D})^T$$

So just set

$$L_1 = L\sqrt{D} \leadsto S = L_1 L_1^T$$

Strang:

" $S = A^T A$  is how a positive-definite symmetric matrix is **put together**."

$S = L_1 L_1^T$  is how you **pull it apart**"

Eg:  $\begin{pmatrix} 2 & 4 & -2 \\ 4 & 9 & -1 \\ -2 & -1 & 14 \end{pmatrix} = L_1 L_1^T$  for

$$L_1 = \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 3 & 1 \end{pmatrix} \begin{pmatrix} \sqrt{2} & 0 & 0 \\ 0 & \sqrt{1} & 0 \\ 0 & 0 & \sqrt{3} \end{pmatrix} = \begin{pmatrix} \sqrt{2} & 0 & 0 \\ 2\sqrt{2} & 1 & 0 \\ -\sqrt{2} & 3 & \sqrt{3} \end{pmatrix}$$

# Quadratic Optimization

This is an important application of the spectral theorem and positive-definiteness. Also,  $SVD + QO + \epsilon\text{-stats} = PCA$ .

It is the simplest case of **quadratic programming**, which is a big subfield of optimization. (So is **least squares**.)

For an example application, see the Wikipedia page for **support-vector machine**, an important tool in machine learning that reduces to a quadratic optimization problem. (There are tons of other applications.)

**Def:** An **optimization problem** means finding extremal values (minimum & maximum) of a function  $f(x_1, \dots, x_n)$  subject to some constraint on  $(x_1, \dots, x_n)$ .

In quadratic optimization, we consider quadratic functions.

**Def:** A **quadratic form** in  $n$  variables is a function  $q(x_1, \dots, x_n) = \text{sum of terms of the form } a_{ij} x_i x_j$

**Eg:**  $q(x_1, x_2) = \frac{5}{2}x_1^2 + \frac{5}{2}x_2^2 - x_1x_2$

**Non-eg:**  $q(x_1, x_2) = x_1^2 + x_2^2 + x_1 + x_2$  is **not** a quadratic form:  $x_1, x_2$  are linear terms.

NB: Thinking of  $x = (x_1, \dots, x_n)$  as a vector,

$$q(cx) = q(cx_1, \dots, cx_n) = \sum a_{ij} (cx_i)(cx_j) \\ = \sum c^2 a_{ij} x_i x_j = c^2 q(x)$$

$$q(cx) = c^2 q(x)$$

In quadratic optimization, the **constraint** on  $x = (x_1, \dots, x_n)$  is usually  $\|x\| = 1$ , i.e.  $x_1^2 + \dots + x_n^2 = 1$ .

### Quadratic Optimization Problem:

Given a quadratic form  $q(x)$ , find the minimum & maximum values of  $q(x)$  subject to  $\|x\| = 1$ .

Eg:  $q(x_1, x_2) = 3x_1^2 - 2x_2^2$

Maximum:

$$q(x_1, x_2) = 3x_1^2 - 2x_2^2 \leq 3x_1^2 + 3x_2^2 \\ = 3(x_1^2 + x_2^2) = 3\|x\|^2 = 3$$

$-2x_2^2 \leq 3x_2^2$

So the maximum value is **3**; it is achieved at  $(x_1, x_2) = \pm(1, 0)$ :  $q(\pm 1, 0) = 3$ .

Minimum:

$$3x_1^2 \geq -2x_2^2$$

$$\begin{aligned} q(x_1, x_2) &= 3x_1^2 - 2x_2^2 \geq -2x_1^2 - 2x_2^2 \\ &= -2(x_1^2 + x_2^2) = -2\|x\|^2 = -2 \end{aligned}$$

So the minimum value is  $-2$ ; it is achieved at  $(x_1, x_2) = \pm(0, 1)$ :  $q(0, \pm 1) = -2$ .

This example is easy because  $q(x_1, x_2) = 3x_1^2 - 2x_2^2$  involves only squares of the coordinates: there is no **cross-term**  $x_1 x_2$ .

**Def:** A quadratic form is **diagonal** if it has the form  $q(x_1, \dots, x_n) = \text{sum of terms of the form } \lambda_i x_i^2$ .

Terms of the form  $a_{ij} x_i x_j$  ( $i \neq j$ ) are **cross-terms**.

**Quadratic Optimization of Diagonal Forms:**

Let  $q(x) = \sum_i \lambda_i x_i^2$ . Order the  $x_i$  so that

$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$ . Then

- The **maximum** value of  $q(x)$  is  $\lambda_1$ .
- The **minimum** value of  $q(x)$  is  $\lambda_n$ .

(subject to  $\|x\|=1$ ).

**NB:** the  $\lambda_i$  could be negative.

**Strategy:** To solve a quadratic optimization problem, we want to **diagonalize** it to get rid of the **cross terms**.

To do this, we use symmetric matrices!

**Fact:** Every quadratic form can be written  
 $q(x) = x^T S x$   
for a symmetric matrix  $S$ .

**Eg:**  $S = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{pmatrix}$

$$\hookrightarrow x^T S x = (x_1 \ x_2 \ x_3) \begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 5 \\ 3 & 5 & 6 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}$$

$$= (x_1 \ x_2 \ x_3) \begin{pmatrix} x_1 + 2x_2 + 3x_3 \\ 2x_1 + 4x_2 + 5x_3 \\ 3x_1 + 5x_2 + 6x_3 \end{pmatrix}$$

$$= x_1^2 + 2x_1x_2 + 3x_1x_3 \\ + 2x_2x_1 + 4x_2^2 + 5x_2x_3 \\ + 3x_3x_1 + 5x_3x_2 + 6x_3^2$$

$$= x_1^2 + 4x_2^2 + 6x_3^2 + 4x_1x_2 + 6x_1x_3 + 10x_2x_3$$

**NB:** The (1,2) and (2,1) entries contribute to the  $x_1x_2$  coefficient.



Given  $q$ , how to get  $S$ ?

The  $x_i^2$  coefficients go on the diagonal, and half of the  $x_i x_j$  coefficient goes in the  $(i,j)$  and  $(j,i)$  entries.

$$q(x_1, x_2, x_3) = a_{11}x_1^2 + a_{22}x_2^2 + a_{33}x_3^2 \\ + a_{12}x_1x_2 + a_{13}x_1x_3 + a_{23}x_2x_3$$

$$\leadsto S = \begin{pmatrix} a_{11} & a_{12}/2 & a_{13}/2 \\ a_{12}/2 & a_{22} & a_{23}/2 \\ a_{13}/2 & a_{23}/2 & a_{33} \end{pmatrix}$$

**NB:**  $q$  is diagonal  $\iff S$  is diagonal: the  $a_{ij}$  are the coefficients of the cross-terms.

$$x^T \begin{pmatrix} \lambda_1 & \dots & 0 \\ & \ddots & \\ 0 & & \lambda_n \end{pmatrix} x = \lambda_1 x_1^2 + \lambda_2 x_2^2 + \dots + \lambda_n x_n^2$$

How does this help quadratic optimization?

Orthogonally diagonalize!

$$q(x) = x^T S x$$

Find a diagonal matrix  $D$  and orthogonal matrix  $Q$  such that  $S = Q D Q^T$

$$\leadsto q(x) = x^T Q D Q^T x$$

Let  $x = Qy$  : this is a change of variables

$$q(x) = q(Qy) = (Qy)^T Q D Q^T (Qy) \\ = y^T \overset{I_n}{Q^T Q} \overset{I_n}{D} Q^T Q y = y^T D y$$

This is now diagonal!

NB:  $Q$  is orthogonal  $\Rightarrow \|x\| = \|Qy\| = \|y\|$

$$\text{So } \|x\|=1 \Leftrightarrow \|y\|=1$$

Eg: Find the minimum & maximum of

$$q(x_1, x_2) = \frac{1}{2}x_1^2 + \frac{1}{2}x_2^2 - 5x_1x_2 \quad \leftarrow \text{cross term} \quad \text{☹️}$$

subject to  $\|x\|=1$ .

$$q(x) = x^T \begin{pmatrix} 1/2 & -5/2 \\ -5/2 & 1/2 \end{pmatrix} x \rightsquigarrow S = \frac{1}{2} \begin{pmatrix} 1 & -5 \\ -5 & 1 \end{pmatrix}$$

Orthogonally diagonalize:  $S = Q D Q^T$  for

$$Q = \frac{1}{\sqrt{2}} \begin{pmatrix} -1 & 1 \\ 1 & 1 \end{pmatrix} \quad D = \begin{pmatrix} 3 & 0 \\ 0 & -2 \end{pmatrix}$$

Set  $x = Qy$ :

$$\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} -1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} -y_1 + y_2 \\ y_1 + y_2 \end{pmatrix}$$

$$\begin{cases} x_1 = \frac{1}{\sqrt{2}}(-y_1 + y_2) \\ x_2 = \frac{1}{\sqrt{2}}(y_1 + y_2) \end{cases} \quad \text{is a linear change of variables}$$

$$\text{Then } q(x) = y^T \begin{pmatrix} 3 & 0 \\ 0 & -2 \end{pmatrix} y = 3y_1^2 - 2y_2^2.$$

Check:

$$\begin{aligned} q(x) &= q\left(\frac{1}{\sqrt{2}}(-y_1+y_2), \frac{1}{\sqrt{2}}(y_1+y_2)\right) \\ &= \frac{1}{2} \cdot \frac{1}{2}(-y_1+y_2)^2 + \frac{1}{2} \cdot \frac{1}{2}(y_1+y_2)^2 - 5 \cdot \frac{1}{2}(-y_1+y_2)(y_1+y_2) \\ &= \frac{1}{4}y_1^2 + \frac{1}{4}y_2^2 - \cancel{\frac{1}{2}y_1y_2} + \frac{1}{4}y_1^2 + \frac{1}{4}y_2^2 + \cancel{\frac{1}{2}y_1y_2} \\ &\quad + \frac{5}{2}y_1^2 - \frac{5}{2}y_2^2 \\ &= \left(\frac{1}{4} + \frac{1}{4} + \frac{5}{2}\right)y_1^2 + \left(\frac{1}{4} + \frac{1}{4} - \frac{5}{2}\right)y_2^2 = 3y_1^2 - 2y_2^2 \end{aligned}$$

The **maximum value** of  $q$  subject to  $\|x\| = \|y\| = 1$  is **3**, achieved at

$$y = (\pm 1, 0) \leadsto x = Qy = \pm \frac{1}{\sqrt{2}} \begin{pmatrix} -1 \\ 1 \end{pmatrix}$$

The **minimum value** of  $q$  subject to  $\|x\| = \|y\| = 1$  is **-2**, achieved at

$$y = (0, \pm 1) \leadsto x = Qy = \pm \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

**NB:** The minimum value is the smallest diagonal entry of  $D \leadsto$  **smallest eigenvalue**.

$Q \begin{pmatrix} 0 \\ \pm 1 \end{pmatrix}$  is  $\pm$  the last column of  $Q$

$\leadsto$  is a **unit eigenvector** for that eigenvalue.

Likewise for the largest eigenvalue.

## Quadratic Optimization:

To find the minimum/maximum of a quadratic form  $q(x)$  subject to  $\|x\|=1$ :

(1) Write  $q(x) = x^T S x$  for a symmetric matrix  $S$

(2) Orthogonally diagonalize  $S = Q D Q^T$  for

$$Q = \underbrace{\begin{pmatrix} | & & | \\ u_1 & \dots & u_n \\ | & & | \end{pmatrix}}_{\text{eigenvectors}} \quad D = \underbrace{\begin{pmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_n \end{pmatrix}}_{\text{eigenvalues}}$$

Order the eigenvalues so  $\lambda_1 \geq \dots \geq \lambda_n$

(3) The maximum value of  $q(x)$  is the largest eigenvalue  $\lambda_1$ .

It is achieved for  $x = \text{any unit } \lambda_1\text{-eigenvector}$

The minimum value of  $q(x)$  is the smallest eigenvalue  $\lambda_n$ .

It is achieved for  $x = \text{any unit } \lambda_n\text{-eigenvector}$ .

**NB:** If  $\dim(\lambda_i) = 1$  then the only unit  $\lambda_i$ -eigenvectors are  $\pm u_i$ . (only 2 unit vectors are on any line)

**NB:**  $x = Qy$  diagonalizes  $q$ :  
 $q(x) = \lambda_1 y_1^2 + \dots + \lambda_n y_n^2$

